

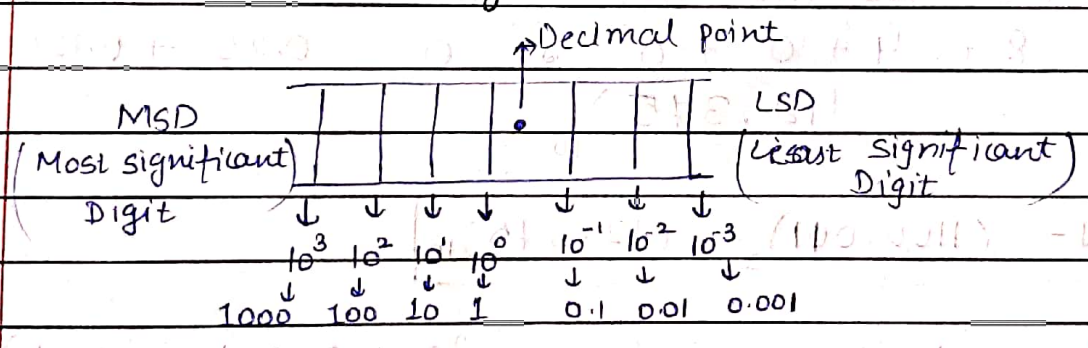
Unit: 1 " Boolean Algebra And Combinational Circuits "

Number System -

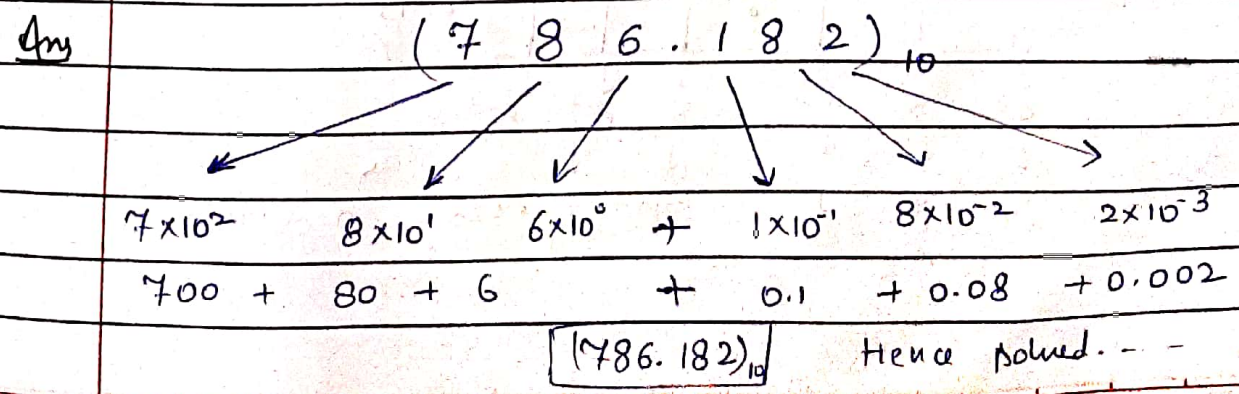
It is a set of rules and symbols used to represent numbers. It's knowledge is very essential because the design and organisation of computer is dependent upon the number system. It is of 4 types -

- | | | | |
|-----|-------------|----|-------------------------|
| (1) | Binary | 2 | } Base (b) or Radix (r) |
| (2) | Octal | 8 | |
| (3) | Decimal | 10 | |
| (4) | Hexadecimal | 16 | |

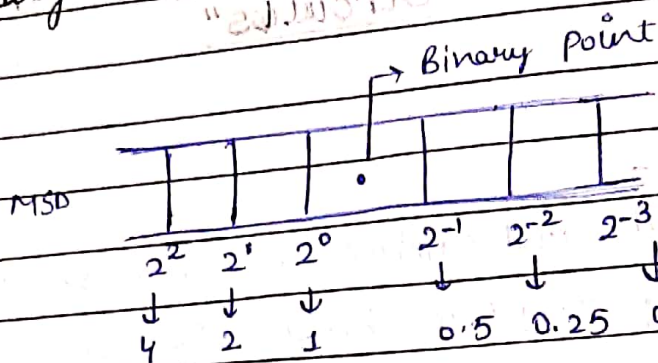
Decimal Number system



Ques Represent decimal number 786.182 in terms of power of 10.



• Binary Number System (0, 1) 1: 3330



Ques Represent the given number in terms of power of 10 and find its decimal equivalent.

$$N = (1100.011)_2$$

Ans

$$(1100.011)_2$$

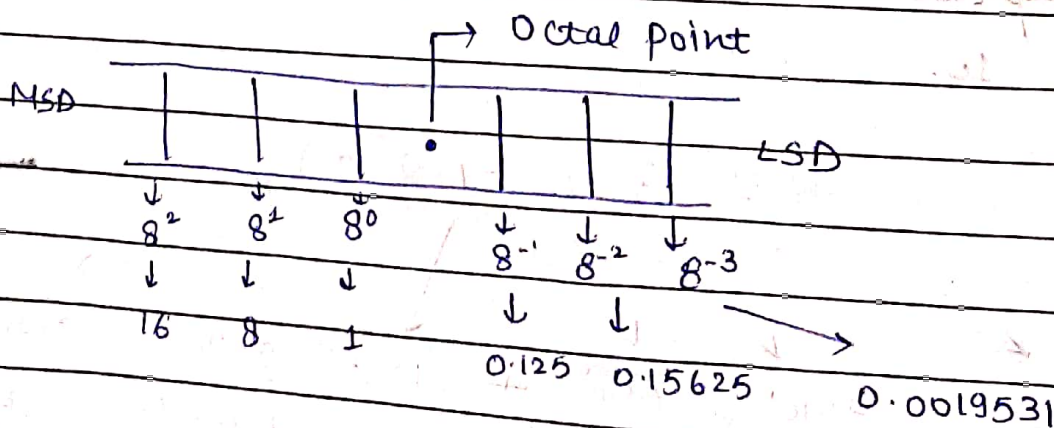
$$1 \times 2^3 + 1 \times 2^2 + 0 \times 2^1 + 0 \times 2^0 + 0 \times 2^{-1} + 1 \times 2^{-2} + 1 \times 2^{-3}$$

$$8 + 4 + 0 + 0 + 0 + 0.25 + 0.125$$

$$(12.375)_{10}$$

$$N = (1100.011)_2 = (12.375)_{10}$$

• Octal Number System (0, 1, 2, 3, 4, 5, 6, 7)



• Conversion from hexadecimal to decimal

$$(2B.48)_{16} \rightarrow (?)_{10}$$

$$(2 \quad B \quad . \quad 4 \quad 8)_{16}$$

$$\downarrow \quad \downarrow \quad \quad \downarrow \quad \downarrow$$

$$2 \times 16^1 \quad 11 \times 16^0 \quad + \quad 4 \times 16^{-1} \quad 8 \times 16^{-2}$$

$$32 + 11 \quad + \quad 0.25 + 0.03125$$

$$43 + 0.28125$$

$$(43.28125)_{10}$$

$$(2B.48)_{16} = (43.28125)_{10}$$

• Conversion from decimal to binary

$$(37.65625)_{10} \rightarrow (?)_2$$

2	37	
2	18	-1
2	9	-0
2	4	-1
2	2	-0
	1	-0

$$0.65625$$

$$\times 2$$

$$\underline{1.31250}$$

↓

1

$$0.625000$$

$$\times 2$$

$$\underline{1.250000}$$

↓

1

$$0.31250$$

$$\times 2$$

$$\underline{0.62500}$$

↓

0

$$0.250000$$

$$\times 2$$

$$\underline{0.500000}$$

↓

0

$$\underline{100101}_2$$

$$0.500000$$

$$\times 2$$

$$\underline{1.000000}$$

↓

1

$$(37.65625)_{10} = (100101.10101)_2$$

$$(10101)_2$$



Conversion from decimal to octal

$(34.45)_{10} \rightarrow (?)_8$

8	34	0.45	0.60	0.80	0.40
8	4 - 2	$\times 8$	$\times 8$	$\times 8$	$\times 8$
	0 - 4	$\underline{3.60}$	$\underline{4.80}$	$\underline{6.40}$	$\underline{3.20}$
		\downarrow	\downarrow	\downarrow	\downarrow
		3	4	6	3

$(42)_8$

0.20
$\times 8$
$\underline{1.60}$
\downarrow
1

$(0.34631)_8$

$(34.45)_{10} = (42.34631)_8$

Conversion from decimal to hexadecimal

$(7825.760)_{10} \rightarrow (?)_{16}$

16	7825	0.760	0.16
16	489 - 1	$\underline{16}$	$\underline{16}$
16	30 - 9	$\underline{12.16}$	$\underline{2.56}$
	1 - 14	\downarrow	\downarrow
		12	2

$(1E91)_{16}$

(14 = E)

0.56
$\underline{16}$
$\underline{8.96}$
\downarrow
8
0.96
$\underline{16}$
$\underline{15.36}$
\downarrow
15

$(C28F)_{16}$



$$(7825.760)_{10} = (1E91.C28F)_{16}$$

Conversion from binary to octal

(1.) $(10110)_2 = (?)_8$

10110

$\begin{array}{c} \text{10110} \\ \downarrow \quad \downarrow \\ 2 \quad 6 \end{array}$

$$(10110)_2 = (26)_8$$

(2.) $(11100110.10111101)_2 = (?)_8$

$\begin{array}{cccccc} \text{011} & \text{100} & \text{110} & . & \text{101} & \text{111} & \text{101} \\ \downarrow & \downarrow & \downarrow & & \downarrow & \downarrow & \downarrow \\ 3 & 4 & 6 & . & 5 & 7 & 5 \end{array}$

$$(11100110.10111101)_2 = (346.575)_8$$

Conversion from binary to Hexadecimal

$$(10110110.10111101)_2 \rightarrow (?)_{16}$$

$\begin{array}{cccccc} \text{1011} & \text{0110} & . & \text{1011} & \text{1100} & \text{1000} \\ \downarrow & \downarrow & & \downarrow & \downarrow & \downarrow \\ 11 & 6 & . & 11 & 12 & 8 \end{array}$

$$(B6.BC8)_{16}$$

$$(10110110.10111101)_2 = (B6.BC8)_{16}$$

Conversion from Octal to Hexadecimal

$(536.21)_8 = (?)_{16}$

$$\begin{array}{cccccc} (5 & 3 & 6 & . & 2 & 1)_8 \\ \downarrow & \downarrow & \downarrow & & \downarrow & \downarrow \\ 101 & 011 & 110 & . & 010 & 001 \end{array}$$

$(101011110.010001)_2$ (octal to binary)

$$\begin{array}{cccccc} 0001 & 0101 & 1110 & . & 0100 & 0100 \\ \hline 1 & 5 & E & . & 4 & 4 \end{array}$$

$(536.21)_8 = (15E.44)_{16}$

Conversion from hexadecimal to octal

$(3CFB.2E)_{16} = (?)_8$

$(3 \ C \ F \ B . 2 \ E)_{16}$

$$\begin{array}{cccccc} (3 & 12 & 15 & 11 & . & 2 & 14)_{16} \\ \downarrow & \downarrow & \downarrow & \downarrow & & \downarrow & \downarrow \end{array}$$

$(0011 \ 1100 \ 1111 \ 1011 . 0010 \ 1110)_2$ (hexadecimal to binary)

$$\begin{array}{cccccc} (000 & 011 & 110 & 011 & 111 & 011 & . & 001 & 011 & 100)_2 \\ \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & & \downarrow & \downarrow & \downarrow \\ (0 & 3 & 6 & 3 & 7 & 3 & . & 1 & 3 & 4)_8 \end{array}$$

$(3CFB.2E)_{16} = (36373.134)_8$



Ques Calculate the value of base X if:

(i) $(225)_x = (341)_8$

(ii) $(100)_x = (61)_8$

Ans (i) $(341)_8$

$$3 \times 8^2 + 4 \times 8^1 + 1 \times 8^0$$

$$192 + 32 + 1$$

$$(225)_{10}$$

A/c to question,

$$(225)_{10} = (225)_x$$

$$\boxed{X = 10}$$

(ii) $(100)_x = (61)_8$

$$6 \times 8^1 + 1 \times 8^0$$

$$48 + 1 = 49$$

$$= 1 \times X^2 + 0 \times X^1 + 0 \times X^0$$

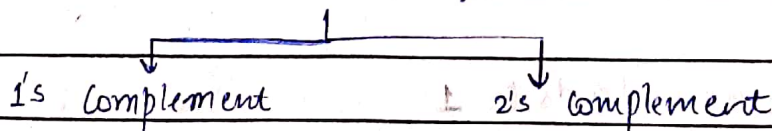
$$= X^2$$

$$X = \sqrt{49}$$

$$\boxed{X = 7}$$

Complement -

Complement of number



(1.) 1's Complement of $(10010)_2$

$$\begin{array}{r} (10010)_2 \\ \downarrow \downarrow \downarrow \downarrow \\ (01101)_2 \end{array}$$

(2.) 2's Complement of $(1010)_2$

$$\begin{array}{r} (1010)_2 \\ \downarrow \downarrow \downarrow \\ (0101)_2 \leftarrow \text{1's complement} \\ +1 \\ \hline (0110)_2 \leftarrow \text{2's complement} \end{array}$$

Unsigned and Signed Binary Numbers -

- The numbers without any positive and negative sign are known as unsigned binary numbers.
- The numbers which are represented by the sign magnitude format are known as signed binary numbers.

Negative (-)

(-) sign tends to 1

example

-6

(-) 6

1 110

= 1110



DATE

positive (+)

(+ sign tends to 0

example +6

(+ sign tends to 0

0 110 = 0110

Decimal Number

Signed Magnitude form

+7

0111

+6

0110

+5

0101

+4

0100

+3

0011

+2

0010

+1

0001

+0

0000

-0

1000

-1

1001

-2

1010

-3

1011

-4

1100

-5

1101

-6

1110

-7

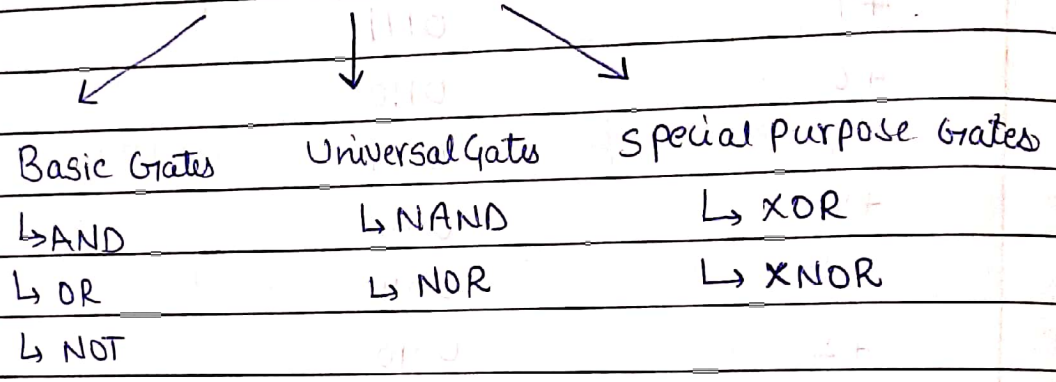
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Signed magnitude form is not possible for 4-bit representation i.e., 8 and more than 8 decimal number.

Logic Gates -

It is an electronic circuit which makes the logical decisions. They have one or more inputs and only one output. Most common logic gates used are AND, OR, NOT, NAND, NOR, XOR, XNOR gates.

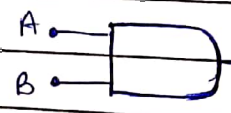
LOGIC GATES



• AND Gate

It gives logical multiplication output. It has two or more inputs and single output.

Mathematically expression, $Y = A \cdot B$



Inputs		Outputs
A	B	Y
0	0	0
0	1	0
1	0	0
1	1	1

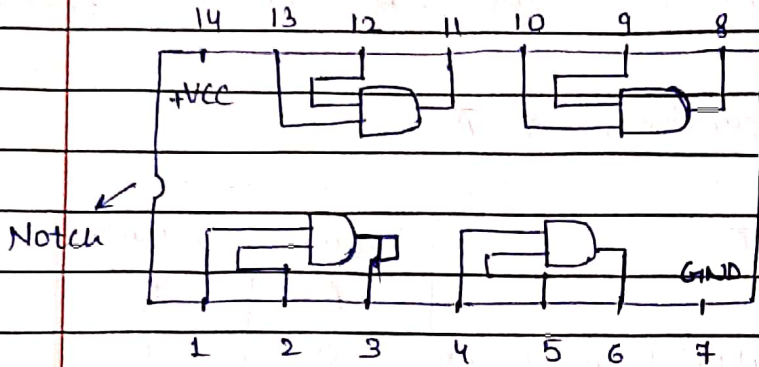
The output of AND Gate is low even if one input is low.



Pin diagram of two input AND gate

~~IC 4708~~

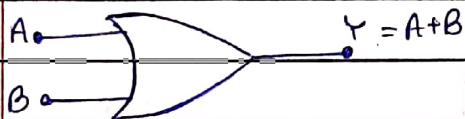
IC 7408



OR Gate

It performs logical addition. It has two or more inputs and single output.

Mathematically expression, $Y = A + B$

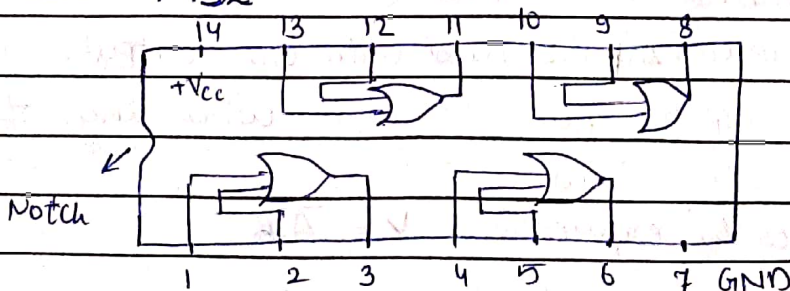


Inputs		Outputs
A	B	Y
0	0	0
1	0	1
0	1	1
1	1	1

Output of OR gate is high only when any one of its input is high.

Pin diagram of two inputs OR gate.

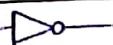
IC 7432



NOT Gate

It performs the basic logical function called complementation. NOT gate is also called inverter. It has one input and one output. It converts one logic gate level into opposite logic level.

Mathematically expression, $Y = \bar{A}$

A  $Y = \bar{A}$

Inputs		Outputs
A		Y
0		1
1		0

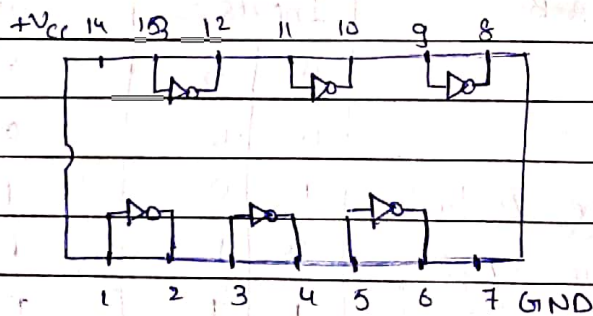
When a high is applied as an input

a low appears on its output and

when a low is applied as an input a high appears on its output.

Pin diagram of NOT gate

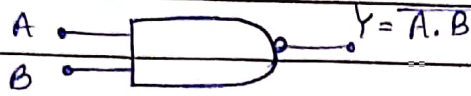
IC 7408



NAND Gate

It is a combination of AND gate and NOT gate. It has two or more inputs and only one output. The output of NAND gate is high when any of the input is low.

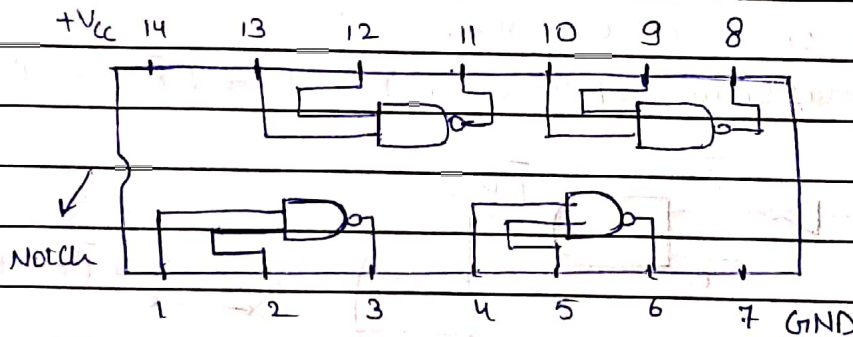
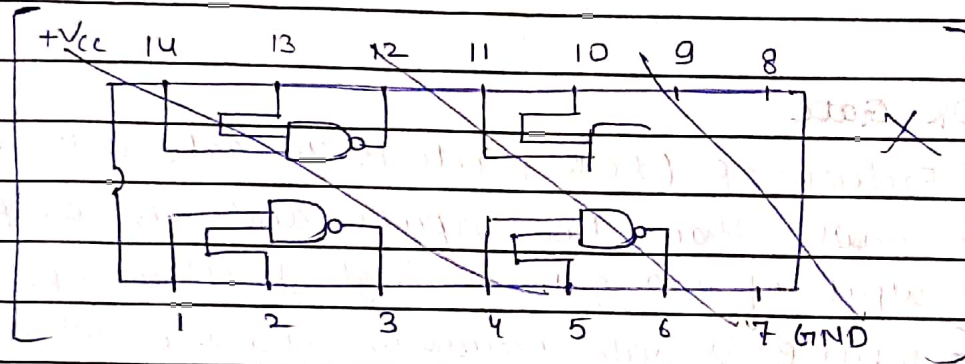
Mathematically expression, $Y = \overline{A \cdot B}$



Inputs		Outputs
A	B	Y
0	0	0
0	1	0
1	0	0
1	1	1

Pin diagram of NAND gate

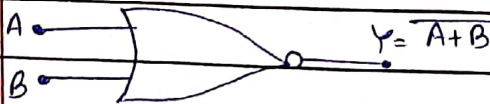
IC 7400



• NOR Gate

It is a combination of OR gate and NOT gate. It has two or more inputs and only one output. The output of NOR gate is high when both inputs are low.

Mathematically expression, $Y = \overline{A+B}$

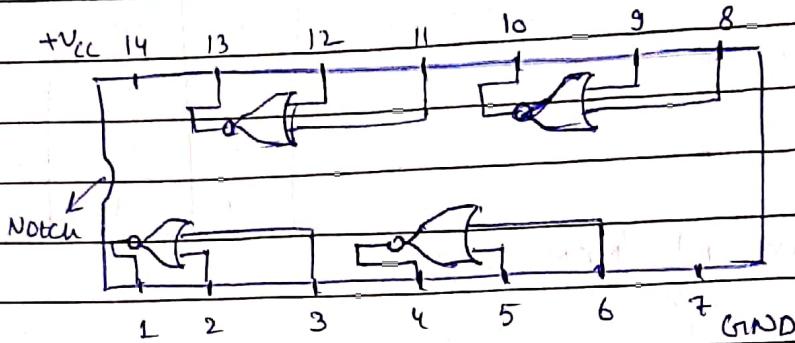


Inputs		Outputs
A	B	Y
0	0	1
0	1	0
1	0	0
1	1	0



Pin diagram of NOR Gate

IC 7402

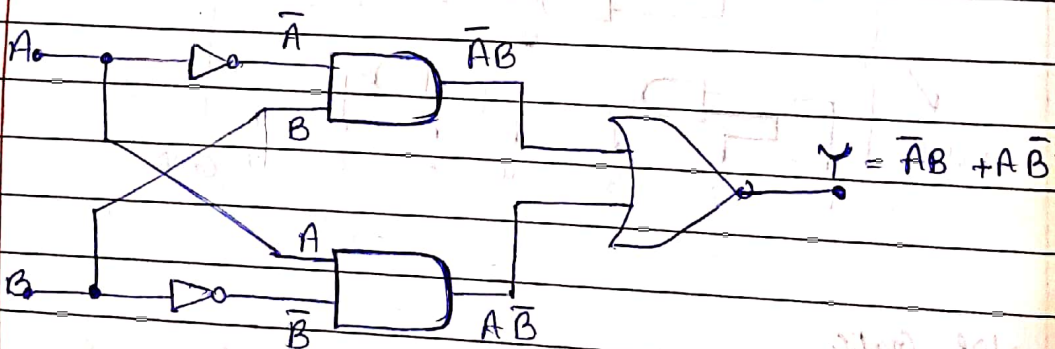


XOR Gate

An Exclusive OR (XOR) gate is the gate with two or more than two inputs and one output.

The output of XOR gate is high if either input A or input B is high exclusively and low when both are '1' or '0' at a time.

Mathematically expression, $Y = \bar{A}B + A\bar{B} = A \oplus B$

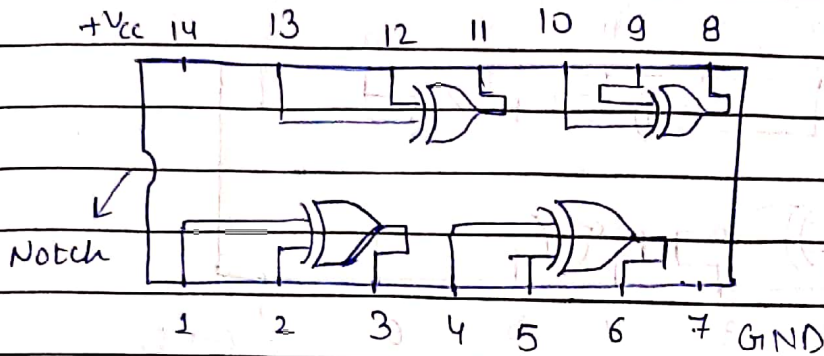


Inputs		Outputs
A	B	Y
0	0	0
0	1	1
1	0	1
1	1	0



Pin diagram of XOR gate

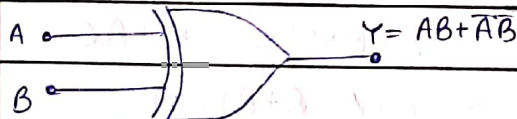
IC 7486



XNOR Gate

An exclusive NOR (XNOR) gate is the gate with two or more than two inputs and one output. The output of XNOR gate is high if both input A and input B are high exclusively or same as input A and input B are low exclusively and low when both are different.

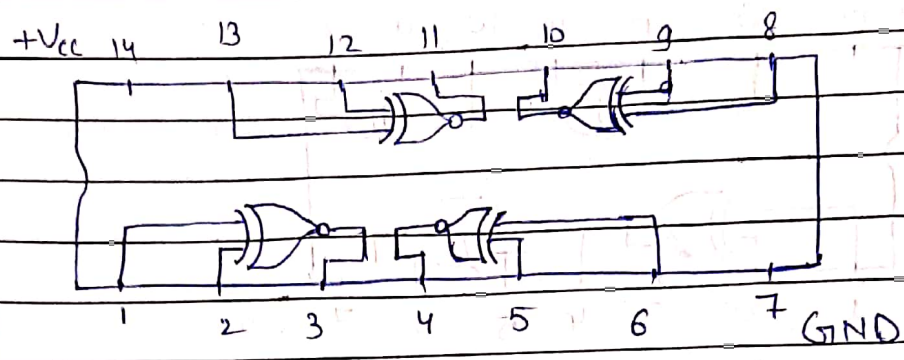
Mathematically expression, $Y = A \oplus B = AB + \bar{A}\bar{B}$



Inputs		Outputs
A	B	Y
0	0	1
0	1	0
1	0	0
1	1	1

$$\begin{aligned}
 \overline{A \oplus B} &= \overline{AB + \bar{A}\bar{B}} \\
 &= \overline{(\bar{A}\bar{B}) + (\bar{A}B)} \\
 &= \overline{(\bar{A} + \bar{B})(\bar{A}B)} \\
 &= \overline{(\bar{A} + B)(A + B)} \\
 &= \bar{A}\bar{A} + \bar{A}\bar{B} + AB + \bar{B}B \\
 &= 0 + \bar{A}\bar{B} + AB + 0 \\
 &= \bar{A}\bar{B} + AB \\
 &= A \oplus B
 \end{aligned}$$

Pin diagram of XOR gate
IC 74266



Theorems of Boolean Algebra -

	$A \cdot 0 = 0$	$A + 0 = A$	
AND laws	$A \cdot 1 = A$	$A + 1 = 1$	OR laws
	$A \cdot A = A$	$A + A = A$	
	$A \cdot \bar{A} = 0$	$A + \bar{A} = 1$	
Commutative law	$A + B = B + A$	$A + B = B + A$	
Associative law	$A \cdot (B \cdot C) = (A \cdot B) \cdot C$	$A + (B + C) = (A + B) + C$	
Distributive laws	$A + (B \cdot C) = (A + B) \cdot (A + C)$	$A(B + C) = AB + AC$	
Absorption law	$A + (A \cdot B) = A$	$A(A + B) = A$	
Double inversion law	$\overline{\bar{A}} = A$		
→ Consensus theorem	$AB + \bar{A}C + BC = AB + \bar{A}C$		
→ De-Morgan's theorem	$\overline{A+B} = \bar{A} \cdot \bar{B}$	$\overline{A \cdot B} = A + \bar{B}$	

De-Morgan's theorem -

(a) $\overline{A+B} = \bar{A} \cdot \bar{B}$



It states that $(\bar{A} \cdot \bar{B})$ is the complement of $A+B$ #

∴ We have to prove that $(A+B) + \bar{A} \cdot \bar{B} = 1$

$$(A+B) + \bar{A} \cdot \bar{B}$$

$$[(A+B) + \bar{A}] [(A+B) + \bar{B}]$$

by distributive law

$$[(A+\bar{A}) + B] [(B+\bar{B}) + A]$$

by Associative law

$$[1+B] [1+A]$$

$$[1] [1]$$

$$1$$

$$\therefore (A+B) + \bar{A} \cdot \bar{B} = 1$$

(b) $\overline{A \cdot B} = \bar{A} + \bar{B}$

It states that $(\bar{A} + \bar{B})$ is the complement of $A \cdot B$

∴ We have to prove that $(A+B) \cdot (\bar{A} \cdot \bar{B}) = 0$

$$(A+B) (\bar{A} \cdot \bar{B})$$

$$A(\bar{A} \cdot \bar{B}) + B(\bar{A} \cdot \bar{B})$$

$$A \cdot \bar{A} \cdot \bar{B} + A \cdot \bar{B} \cdot \bar{B}$$

$$0 \cdot \bar{B} + \bar{B} \cdot 0$$

$$0 + 0$$

$$0$$

$$\therefore (A+B) \cdot (\bar{A} \cdot \bar{B}) = 0$$

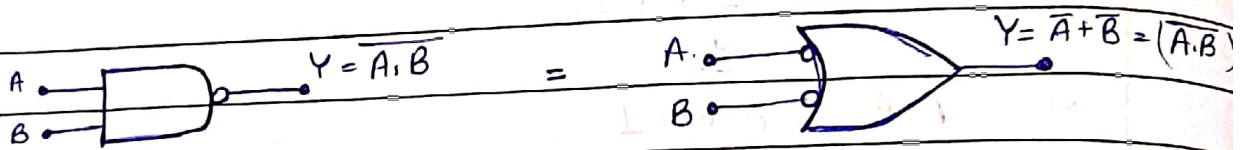
Hence, by using (1) and (2) we concluded that $(\bar{A} \cdot \bar{B})$ is the complement of $(A+B)$

Thus, $\overline{(A+B)} = (\bar{A} \cdot \bar{B})$ and $\overline{(\bar{A} \cdot \bar{B})} = (\bar{A} + \bar{B})$

Universal Gates -

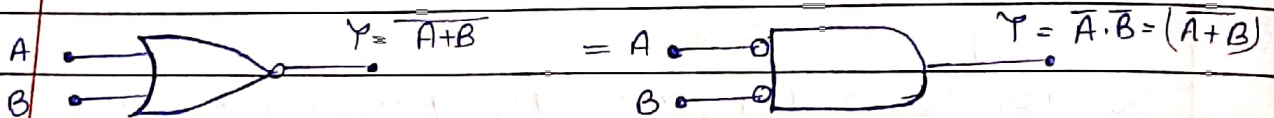
These are those gates which can be used for implementing any gate like AND, OR, XOR, XNOR, NOT or any combination of basic gates.

NAND Gate



NAND gate is also known as bubbled OR gate

NOR Gate

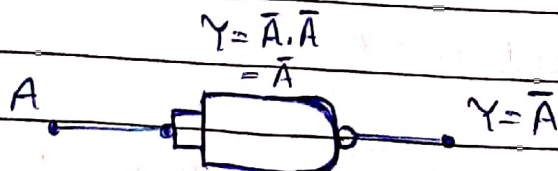


NOR gate is also known as bubbled AND gate

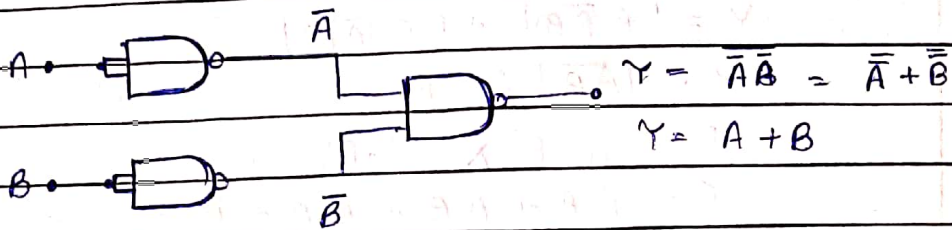
- Universal property of NAND gate

NAND is called universal gate because it can be realize any gate.

→ Realize NOT gate using NAND gate



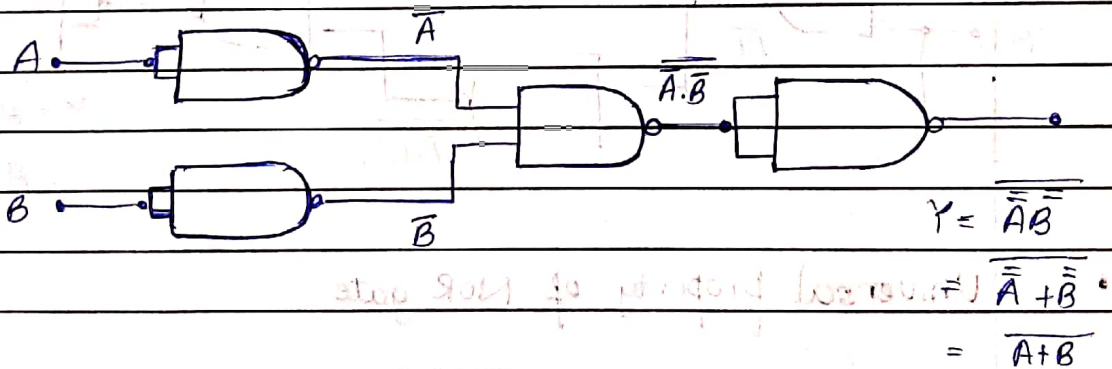
→ Realize OR gate using NAND gates



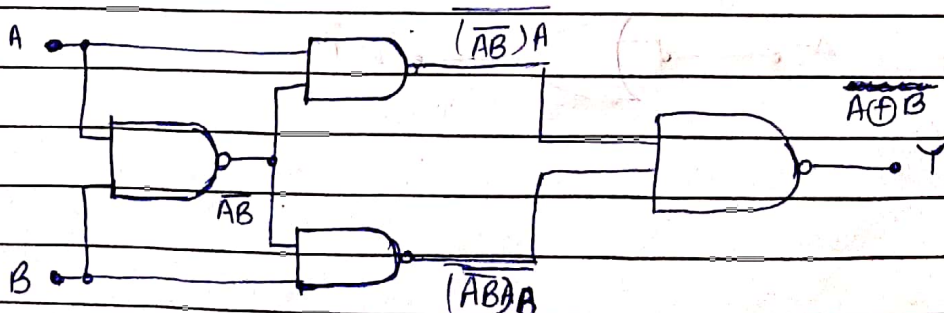
→ Realize AND gate using NAND gates



→ Realize NOR gate using NAND gates



→ Realize a two input EXOR gates using NAND gates



$$Y = \overline{[A \cdot (\overline{A \cdot B})]} [B \cdot (\overline{A \cdot B})]$$

$$Y = \overline{[A \cdot (\overline{A \cdot B})]} + [B \cdot (\overline{A \cdot B})]$$

$$Y = [A \cdot (\overline{A \cdot B})] + [B \cdot (\overline{A \cdot B})]$$

$$Y = [A \cdot \overline{B}] (A+B)$$

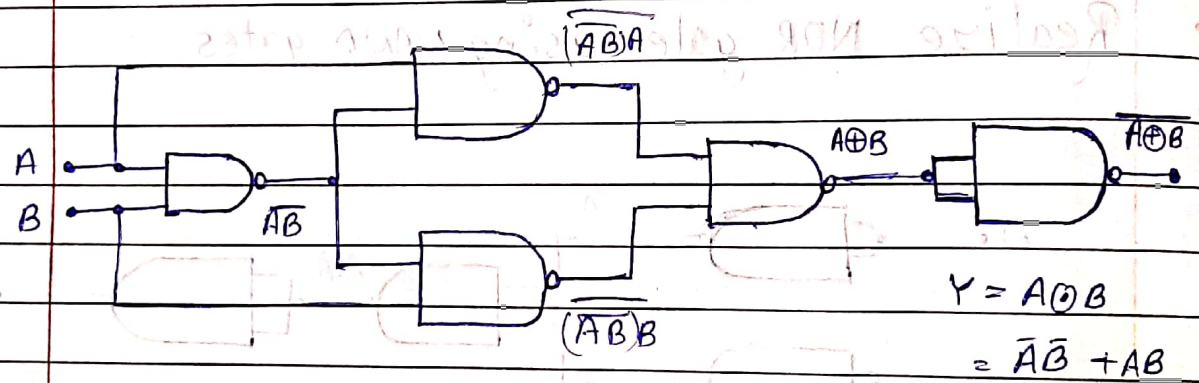
$$Y = (\overline{A+B}) \cdot (A+B)$$

$$Y = \overline{A} \cdot A + \overline{A} \cdot B + \overline{B} \cdot A + \overline{B} \cdot B$$

$$Y = \overline{A} \cdot B + \overline{B} \cdot A = A \oplus B$$

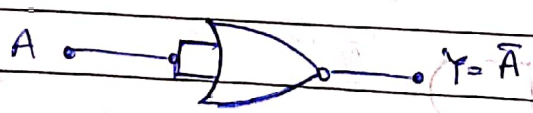
• ~~Universal property of NOR gate~~

→ Realize XNOR gate using NAND gate



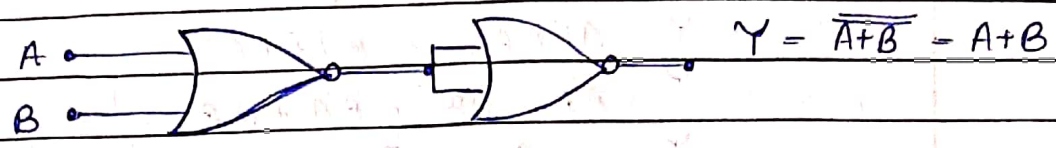
• Universal property of NOR gate

→ Realize NOT gate using NOR gate

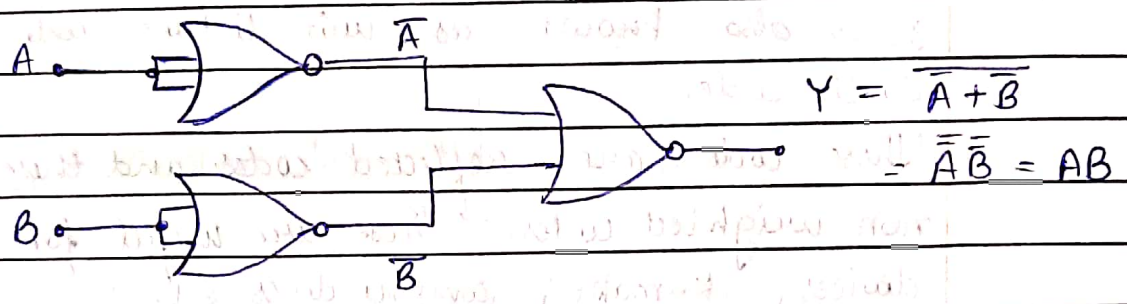


OR

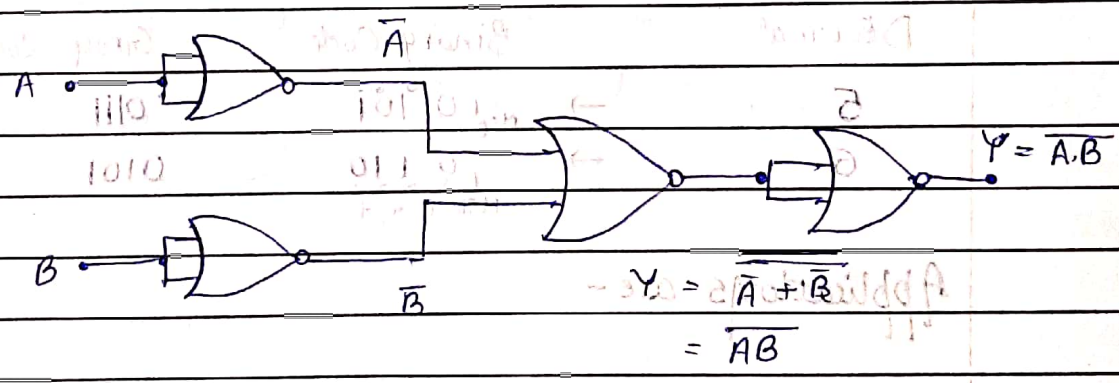
→ Realize ~~AND~~ gate using NOR gate



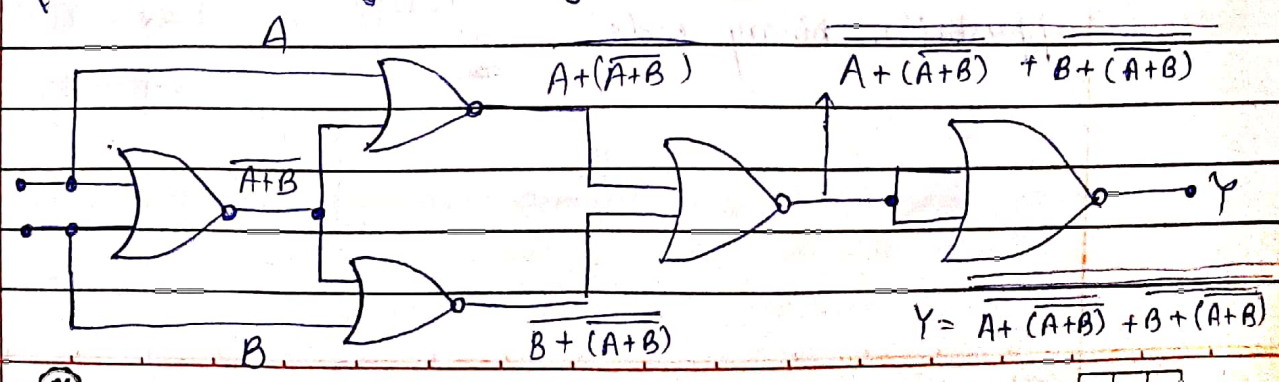
→ Realize AND gate using NOR gate



→ Realize NAND gate using NOR gate



→ Realize EXOR gate using NOR gate



$$Y = \overline{A + (A+B)} + B + (A+B)$$

$$Y = \overline{A} \cdot \overline{[A+B]} + B \cdot \overline{[A+B]}$$

$$Y = \overline{A}[A+B] + B[A+B]$$

$$Y = \overline{A}A + \overline{A}B + BA + BB$$

$$Y = \overline{A}B + AB = A \oplus B$$

#

Gray Code

It is also known as unit distance code and cyclic code.

These codes are reflected codes and these are non-weighted codes. These are useful for input-output devices, k-maps, compact discs ect.

Example -

Decimal		Binary Code	Gray Code
5	→	$\begin{array}{c} \text{MSB} \leftarrow 0 \ 1 \ 0 \ 1 \\ \leftarrow \leftarrow \leftarrow \leftarrow \end{array}$	0111
6	→	$\begin{array}{c} 0 \ 1 \ 1 \ 0 \\ \leftarrow \leftarrow \leftarrow \leftarrow \\ \text{MSB} \leftarrow \leftarrow \leftarrow \leftarrow \end{array}$	0101

Applications are -

- It is mainly used in shaft position encoders.
- These are used in the optical discs to produce an appropriate binary code. (CD's)



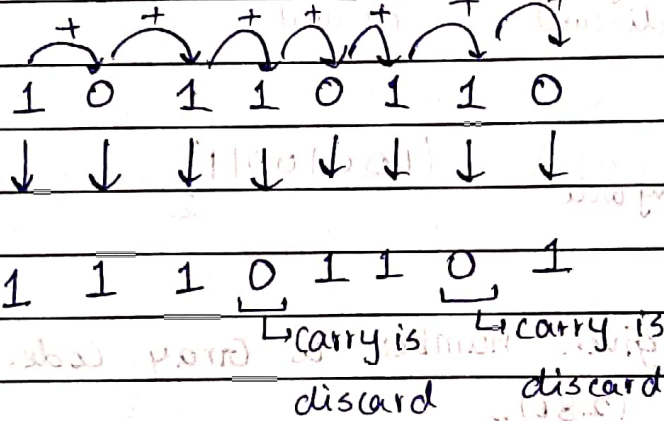
Binary to Gray code conversion -

- The MSB of the binary number remains same in Gray code.
- Move from MSB to LSB and add each adjacent pair of binary digits, if carry comes neglect or discard the carry and write the sum as next Gray code digit.

Ques

Convert $(10110110)_2$ to Gray code.

Ans



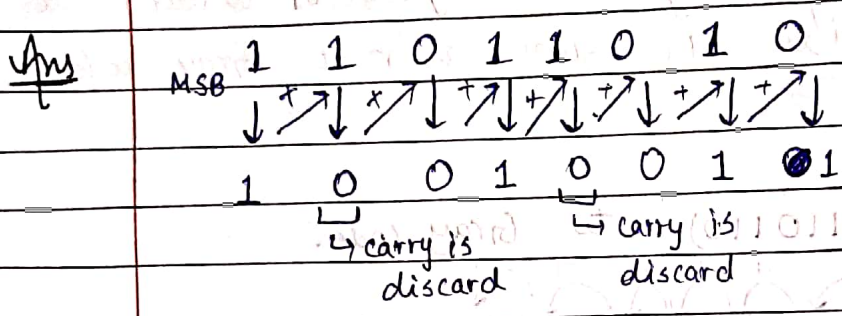
$$(10110110)_2 = (11101101)_{\text{Gray code}}$$

Gray code to Binary Conversion

- The MSB of the Gray code remains same in binary number.
- Add each binary digit generated to the next significant

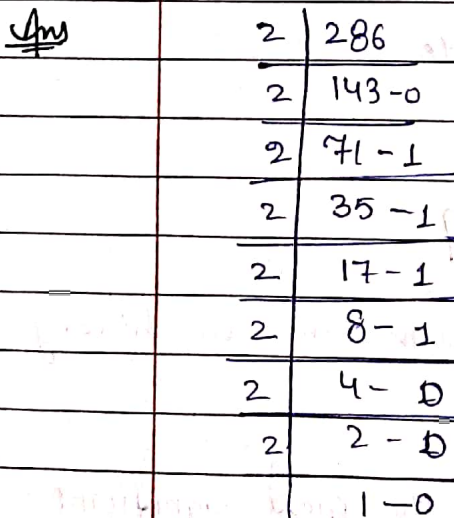
bit of the gray-code, if carry comes, discard the carry and write the sum as next binary bit.

Ques Convert a given Gray Code to binary
 (11011010) Gray code



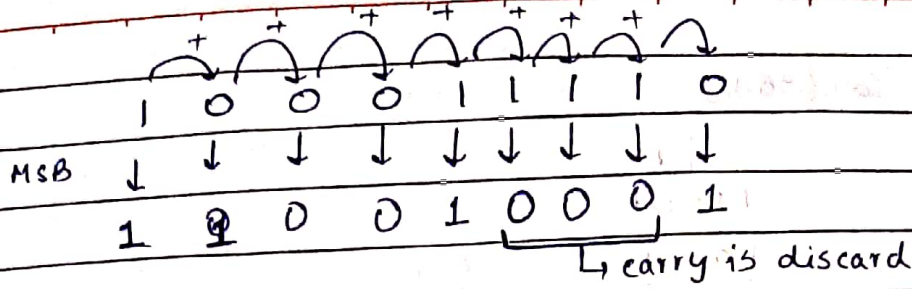
(11011010) Gray code = $(10010011)_2$

Ques Convert a given number to Gray Code.
 $(286)_{10}$



$(286)_{10} = (10001110)_2$





$(100011110)_2 = (110010001)_{\text{Gray code}}$

Complement arithmetic -

Ques Do the following

- (i) $(25)_{10} - (13)_{10}$ in 2's complement arithmetic
- (ii) $(13)_{10} - (25)_{10}$ in 2's complement arithmetic

Ans (i) $(25)_{10} + (-13)_{10}$

2	25	2	13
2	12 - 1	2	6 - 1
2	6 - 0	2	3 - 0
2	3 - 0		1 - 1
	1 - 1		

$(13)_{10} = (1101)_2$

$(25)_{10} = (11001)_2$

→ 2's complement of 13,

1's complement = 10010

1 tends to negative sign

2's complement = 10011

Add it to $(25)_{10}$

$$\begin{array}{r}
 11001 \\
 10011 \\
 \hline
 \text{Carry} \leftarrow \textcircled{1} \quad 01100 \\
 \hline
 \text{discard}
 \end{array}$$

$$(11001)_2 - (01101)_2 = (01100)_2$$

$$(25)_{10} - (13)_{10} = -(12)_{10}$$

(ii)

$$(13)_{10} + (-25)_{10}$$

2	13	25
2	6-1	13-1
2	3-0	6-0
	1-1	3-0
		1-1

$$(13)_{10} = (1101)_2$$

$$(25)_{10} = (11001)_2$$

2's complement of $(25)_{10}$

1's complement $\rightarrow 00110$

+1

2's complement $\rightarrow 00101$

Add it to $(13)_{10}$

$$\begin{array}{r}
 01101 \\
 01101 \\
 \hline
 00111 \\
 \hline
 10100
 \end{array}$$



Since there is no carry, obtain 2's complement and assign negative sign,

$$\begin{array}{r} (10100)_2 \\ = 01011 \\ + 1 \\ \hline -01100 \end{array}$$

$$(01101)_2 - (11001)_2 = (-01100)_2$$

$$(13)_{10} - (25)_{10} = (-12)_{10}$$

Ques

Do $-58 - 13$ in 2's complement arithmetic.

Sol

$$\begin{array}{r} 2 \overline{) 58} \\ \underline{2 \ 29 - 0} \\ 2 \ 14 - 1 \\ \underline{2 \ 7 - 0} \\ 2 \ 3 - 1 \\ \underline{1 - 1} \end{array}$$

$$\begin{array}{r} 2 \overline{) 13} \\ \underline{2 \ 6 - 1} \\ 2 \ 3 - 0 \\ \underline{1 - 1} \end{array}$$

$$(+58)_{10} = (111010)$$

for 8 bits

for 8 bits $(+58)_{10} = (00111010)_2$

$$(+13)_{10} = (00001101)_2$$

$$(+58)_{10} = (00111010)_2$$

$$(-13) \text{ 1's complement} = 11110010$$

$$(-58) \text{ 1's complement} = 11000101$$

+ 1

$$2's \text{ complement} = 11110011$$

$$2's \text{ complement} = \underline{11000110}$$



$$(-58) + (-13)$$

$$\begin{array}{r} 11000110 \\ + 11110011 \\ \hline 110111001 \end{array}$$

↓
discard carry

2's complement of $(10111001)_2$

$$\begin{array}{r} 1's \text{ complement} \rightarrow 01000110 \\ + 1 \end{array}$$

$$2's \text{ complement } \underline{01100111}$$

sign bit

$$(-58)_{10} + (-13)_{10} = -(71)_{10}$$

Ques

Subtract the following decimal numbers using 9's complement:

(i) $(23)_{10} - (12)_{10}$

(ii) $(236.34)_{10} - (105.93)_{10}$

(iii) $(17)_{10} - (6)_{10}$

(iv) $(105.23)_{10} - (236.34)_{10}$

Ans

(i) $23 - 12$

The 9's complement of 12 is,

$$\begin{array}{r} 99 \\ - 12 \\ \hline 87 \end{array}$$

$$\begin{array}{r}
 23 \\
 + 87 \\
 \hline
 \textcircled{1} 10 \\
 \swarrow +1 \\
 \hline
 11
 \end{array}$$

← (Add 9's complement of 12)

← (Add end 'around' carry)

← final result

$$\boxed{(23)_{10} - (12)_{10} = (11)_{10}}$$

(ii) 17-6

The 9's complement of 6 is,

$$\begin{array}{r}
 99 \\
 -06 \\
 \hline
 93
 \end{array}$$

← (9's complement of 6)

$$\begin{array}{r}
 17 \\
 + 93 \\
 \hline
 \textcircled{1} 10 \\
 \swarrow +1 \\
 \hline
 11
 \end{array}$$

← (Add 9's complement of 6)

$$\boxed{(17)_{10} - (6)_{10} = (11)_{10}}$$

(iii) $(236.34)_{10} - (105.23)_{10}$

$$\begin{array}{r}
 999.99 \\
 -105.23 \\
 \hline
 894.76
 \end{array}$$

← (9's complement of 105.23)

#

$$\begin{array}{r}
 238.34 \\
 + 894.76 \\
 \hline
 \textcircled{1} 138.10 \\
 \xrightarrow{+1} \\
 \hline
 138.11 \quad \text{Ans.}
 \end{array}$$

(iv) $(105.23)_{10} - (236.34)_{10}$

$$\begin{array}{r}
 999.99 \\
 - 236.34 \\
 \hline
 763.65 \quad \leftarrow \text{(9's complement of } 236.34)
 \end{array}$$

$$\begin{array}{r}
 105.23 \\
 + 763.65 \\
 \hline
 868.88
 \end{array}$$

(v) There is no carry at MSD, thus the result is -ve.

$$\begin{array}{r}
 999.99 \\
 - 868.88 \\
 \hline
 - 131.11 \quad \leftarrow \text{(9's complement of result)} \\
 \text{and assign -ve sign}
 \end{array}$$

Ques Subtract $(18)_{10} - (24)_{10}$ using 9's complement

$$\begin{array}{r}
 99 \\
 - 24 \\
 \hline
 75 \quad \leftarrow \text{(9's complement of } -24)
 \end{array}$$



DATE

$$\begin{array}{r} 18 \\ + 75 \\ \hline 93 \end{array}$$

There is no carry at MSD, thus the result is -ve.

$$\begin{array}{r} 99 \\ - 93 \\ \hline 06 \end{array} \leftarrow (9\text{'s complement of result) and assign -ve sign$$

$$(18)_{10} - (24)_{10} = (-6)_{10}$$

Ques Subtract the following decimal numbers using 10's complement:

(i) $(732.19)_{10} - (213.42)_{10}$

(ii) $(213.42)_{10} - (732.19)_{10}$

Ans (i)
$$\begin{array}{r} 999.99 \\ - 213.42 \\ \hline 786.57 \end{array} \leftarrow (9\text{'s complement of } 213.42)$$

$$\begin{array}{r} 786.57 \\ + 1 \\ \hline 786.58 \end{array} \leftarrow (10\text{'s complement of } 213.42)$$

$$\begin{array}{r}
 732.19 \\
 786.58 \\
 \hline
 \textcircled{1} 518.77 \\
 \hline
 \text{+ discard carry} \\
 \hline
 \hline
 \hline
 \end{array}$$

$$(732.19)_{10} - (213.42)_{10} = (518.77)_{10}$$

(ii)

$$\begin{array}{r}
 999.99 \\
 - 732.19 \\
 \hline
 267.80 \quad \leftarrow \text{(9's complement of } 732.19) \\
 + 1 \\
 \hline
 267.81 \quad \leftarrow \text{(10's complement of } 732.19)
 \end{array}$$

$$\begin{array}{r}
 267.81 \\
 + 213.42 \\
 \hline
 481.23
 \end{array}$$

There is no carry so the result is -ve.

$$\begin{array}{r}
 999.99 \\
 - 481.23 \\
 \hline
 518.76 \\
 + 1 \\
 \hline
 518.77 \quad \leftarrow \text{(10's complement and assign -ve sign)}
 \end{array}$$

$$(213.42)_{10} - (732.19)_{10} = (-518.77)_{10}$$



SOP and POS -

Sum of Product form (SOP)

In this case product terms of a function are added.

In this '0' \rightarrow complemented form = \bar{A}

'1' \rightarrow non-complemented form = A

Sum of product ^{terms} are known as 'Minterms' and denoted by 'm'.

Inputs			SOP terms	Minterms Abbreviations
A	B	C		
0	0	0	$\bar{A}\bar{B}\bar{C}$	m_0
0	0	1	$\bar{A}\bar{B}C$	m_1
0	1	0	$\bar{A}B\bar{C}$	m_2
0	1	1	$\bar{A}BC$	m_3
1	0	0	$A\bar{B}\bar{C}$	m_4
1	0	1	$A\bar{B}C$	m_5
1	1	0	$AB\bar{C}$	m_6
1	1	1	ABC	m_7

Product of sum form (POS)

In this case sum terms of a function are multiplied.

In this '0' \rightarrow non-complemented form = A

'1' \rightarrow complemented form = \bar{A}

Product of sum terms are known as 'Maxterms' and denoted by 'M'.

Inputs			standard - POS terms	Maxterms Abbreviations
A	B	C		
0	0	0	$A+B+C$	M_0
0	0	1	$A+B+\bar{C}$	M_1
0	1	0	$A+\bar{B}+C$	M_2
0	1	1	$A+\bar{B}+\bar{C}$	M_3
1	0	0	$\bar{A}+B+C$	M_4
1	0	1	$\bar{A}+B+\bar{C}$	M_5
1	1	0	$\bar{A}+\bar{B}+C$	M_6
1	1	1	$\bar{A}+\bar{B}+\bar{C}$	M_7

Karnaugh Map -

It is simply a graphical method for representing a boolean function. It is used for the minimization of switching functions but upto 'six' variables.

Types of k-Maps commonly used -

- | | | |
|-----|----------------------|--------------------|
| | | It consists of |
| (1) | Two Variable k-map | $2^2 = 4$ squares |
| (2) | Three variable k-map | $2^3 = 8$ squares |
| (3) | Four variable k-map | $2^4 = 16$ squares |
| (4) | Five variable k-map | $2^5 = 32$ squares |
| (5) | Six variable k-map | $2^6 = 64$ squares |

Ques Minimize the given function using k-map.
 $Y(A,B,C) = \sum m(1,3,5,7)$

Ans

		BC	$\overline{B}\overline{C}$	$\overline{B}C$	BC	BC
A		00	01	11	10	
\overline{A}	0	0	1 ₁	1 ₃	0	
A	1	0	1 ₅	1 ₇	0	

$\therefore Y = C$

Ques Minimize the given function using k-map.
 $Y(A,B,C,D) = \sum m(0,2,5,7,8,10,13,15)$

Ans

		CD	$\overline{C}\overline{D}$	$\overline{C}D$	CD	$C\overline{D}$
AB		00	01	11	10	
$\overline{A}\overline{B}$	00	1 ₀	0	0	1 ₃	1
$\overline{A}B$	01	0	1 ₅	1 ₇	0	
$A\overline{B}$	11	0	1 ₁₃	1 ₁₅	0	
AB	10	1 ₈	0	0	1 ₁₀	1

Group 1

$Y = \text{Group 1} + \text{Group 2} + \text{Group 3}$
 $Y = BD + \overline{B}\overline{D}$

Ques Minimize the given function using k-map
 $Y(A,B,C,D) = \sum m(0,1,2,3,6,8,9,10,11,12,13)$

Ans

		CD	$\overline{C}\overline{D}$	$\overline{C}D$	$C\overline{D}$	CD
AB		00	01	11	10	
$\overline{A}\overline{B}$	00	0	0	0	0	
$\overline{A}B$	01	0	0	0	0	
$A\overline{B}$	11	0	0	0	0	
AB	10	0	0	0	0	

Group 1 B

Group 2 $A + \overline{C} + D$

Group 4 $C + \overline{A}$



$$Y = (B)(A + \bar{C} + D)(\bar{A} + C)$$

Ques Simplify the following logic expression using K-map.

$$Y(A, B, C, D, E) = \sum m(3, 5, 6, 8, 9, 12, 13, 14, 19, 22, 24, 25, 30)$$

Ans

BC	DE				BC	DE			
	00	01	11	10		00	01	11	10
00	0	1	1	2	00	16	17	18	19
01	4	1	5	6	01	20	21	22	23
11	12	1	13	14	11	28	29	30	31
10	8	1	9	10	10	24	25	26	27

A=0 or \bar{A}

A=1 or A

Group 1 = $\bar{B}\bar{C}DE$

Group 2 = CDE

Group 3 = $B\bar{C}\bar{D}$

Group 4 = $\bar{A}CDE$ (\bar{A} include because of left most K-map)

Group 5 = $\bar{A}B\bar{D}$ ()

$$Y = \bar{B}\bar{C}DE + CDE + B\bar{C}\bar{D} + \bar{A}CDE + \bar{A}B\bar{D}$$

Ques Minimize the following function using K-map.

$$Y(A, B, C, D, E) = \pi M(2, 7, 8, 13, 23, 24, 29)$$

BC	DE 00	DE 01	DE 11	DE 10	BC	DE 00	DE 01	DE 11	DE 10	
BE 00	0	1	3	2	1	BE 00	16	17	19	18
B \bar{E} 01	4	5	0	6	2	B \bar{E} 01	20	21	0	23
$\bar{A}\bar{E}$ 11	12	0	13	15	3	$\bar{A}\bar{E}$ 11	24	0	25	31
$\bar{A}E$ 10	0	8	9	11	4	$\bar{A}E$ 10	0	24	25	27

Group 1 = $(A + B + C + \bar{D} + E)$

Group 2 = $(B + \bar{C} + \bar{D} + \bar{E})$

Group 3 = $(\bar{B} + \bar{C} + D + \bar{E})$

Group 4 = $(\bar{B} + \bar{C} + D + E)$

$$Y = (A + B + C + \bar{D} + E)(B + \bar{C} + \bar{D} + \bar{E})(\bar{B} + \bar{C} + D + \bar{E})(\bar{B} + \bar{C} + D + E)$$

Ques Minimize the given expression using K-map.

$$Y(A, B, C, D, E, F) = \sum m(4, 5, 6, 7, 8, 9, 20, 21, 22, 23, 24, 25, 34, 35, 36, 37, 38, 39, 50, 51, 52, 53, 54, 55, 56, 57)$$

A/B	B=0				B=1			
-----	-----	--	--	--	-----	--	--	--

A/B	CD \ EF	00	01	11	10	A/B	CD \ EF	00	01	11	10
A=0	0	1	3	2	16	17	19	18	20	21	23
1	4	5	6	7	24	25	27	26	28	29	30
1	12	13	15	14	34	35	37	36	38	39	40
1	8	9	11	10	44	45	47	46	48	49	50
A=1	32	33	35	34	54	55	57	56	58	59	60
1	20	21	23	22	64	65	67	66	68	69	70
1	44	45	47	46	74	75	77	76	78	79	80
1	40	41	43	42	84	85	87	86	88	89	90

$$Y(A, B, C, D, E, F) = G_1 + G_2 + G_3 + G_4$$

Group 1 = $\bar{C}D$
 (Common for all 4 variable k-map, thus A and B variables are directly eliminated).

Group 2 = $\bar{A}C\bar{D}\bar{E}$

Group 3 = $B C \bar{D} \bar{E}$

Group 4 = $A \bar{C} \bar{E}$

$$Y(A, B, C, D, E, F) = \bar{C}D + \bar{A}C\bar{D}\bar{E} + B C \bar{D} \bar{E} + A \bar{C} \bar{E}$$

Magnitude Comparator

Comparator compares two binary numbers.

Example-

Single bit magnitude comparator,



Inputs		Outputs		
A	B	A > B	A = B	A < B
0	0	0	1	0
0	1	0	0	1
1	0	1	0	0
1	1	0	1	0